

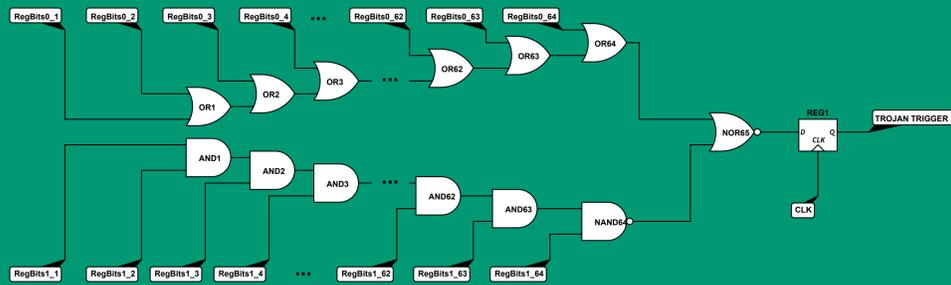
The Design and Implementation of an Open-Source Hardware Trojan for a 64-bit RISC-V CPU Design

1. MOTIVATION

- Hardware Trojan (HT) paradigms on modern, realistic and complex designs are scarce. Development of corresponding open-source HT-testbeds is important for research conducted on defense mechanisms and detection methods for HT.
- The silicon itself can enable attacks that disable or selectively by-pass fundamental security mechanisms (e.g., memory protection) of modern Central Processing Units (CPUs).
- Overwriting data inside the kernel address space from a user process violates address space isolation, a powerful Operating System security mechanism.
- Current HT detection methods pertain a prohibitive economic cost, are very time consuming and can lead to the destruction of the Device Under Test.

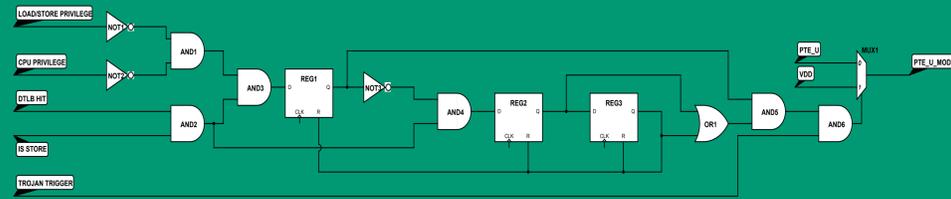
3. DESIGN & IMPLEMENTATION OF ASTRAHAN

Triggering Circuit



The **trigger** circuit targets Ariane's **General-Purpose Registers (GPRs)**, as it monitors them for specific byte values to assert the trigger signal. Triggering circuit's design is **scalable**. For the worst-case scenario of a 128-bit magic value, it requires only **130 logic gates**.

Payload Circuit



The **payload** circuit targets Ariane's **Memory Management Unit (MMU)**. Payload circuit design has a **minimal footprint**, as it requires only **14 logic gates**.

FPGA resource utilization of Astrahan's circuits and Ariane's implementation.

Module	LUTs	FFs	LUT Utilization	FF Utilization
Ariane*	66133	51854	32.45%	12.72%
Register File**	1240	1985	1.87%	3.82%
Trigger Circuit**	26	1	0.00039%	0.000019%
Memory Management Unit**	4271	3388	6.45%	6.53%
Payload Circuit**	4	3	0.0006%	0.000057%

* The utilization percentages referenced here are with respect to the total FPGA resources.
** The utilization percentages referenced here are with respect to the total Ariane Implementation.

5. ASTRAHAN VS A2

Trojan Name	Gate Number	Host Module	Operational Flexibility	Activation Status	ASIC Feasibility	Open-Source Design
A2	1 & 91	OR1200 (Open RISC)	High	Check needed	Feasible	Yes
Astrahan	144***	Ariane (RISC-V)	Very High	Check not needed	Feasible	Yes

- Open-RISC vs RISC-V ISA.
- Equivalent number of gates for purely digital trojan implementation.
- Astrahan's "context switch safe" design provides very high operability during attack time.

- Astrahan's design requires no prior status activation validation during attack time.
- Astrahan's design provides the **scalability** required during a fabrication time attack.
- Astrahan's design is **open-sourced** to the community just like A2's design.

*** The number of gates for the worst-case scenario, where an attacker chooses to use a 128-bit value as the triggering input.

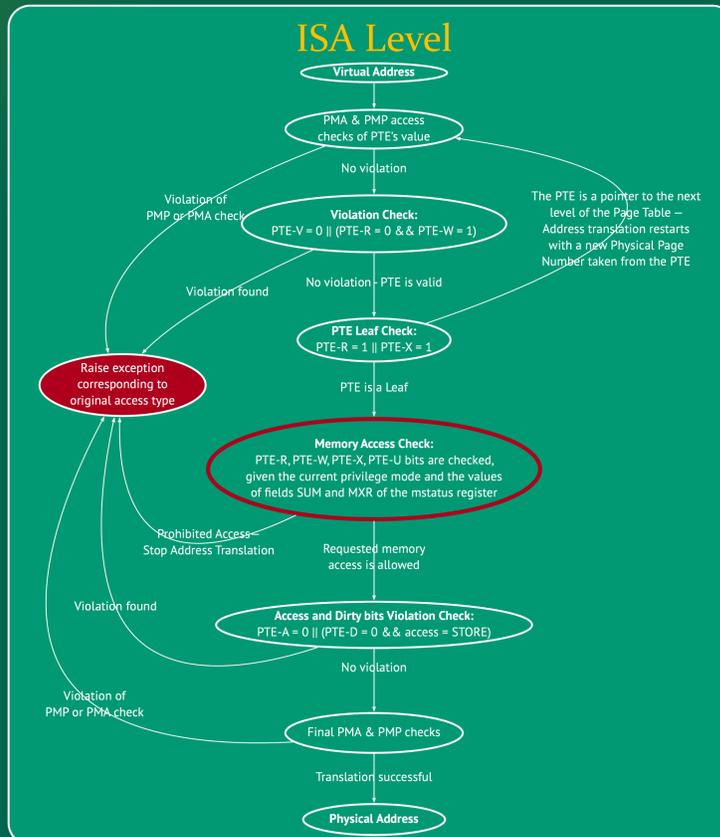
6. THE ASIC SCENARIO

Astrahan's **low-footprint** and **scalable** design alleviates insertion during **fabrication** as:

- The GPRs and the number of bits utilized by the triggering circuit, can be chosen according to the GPRs' placement and the layout's overall spatial options.
- The time sensitive path of the triggering signal is a **multicycle-path**, as a page table walk is initiated every time the user process accesses a kernel space virtual address. Thus, Astrahan's design can cope with demanding layouts where the payload and triggering circuit need to be placed far apart.
- Astrahan's design uses a minimum number of flip-flops, thus putting minimum stress on the local clock networks and requiring less space for placement.

The overhead of the trojan gates on the original circuit paths can be calculated by exporting in test-benches only the layout part of concern for accurate RC parasitic extractions. This enables optimal choices for the gate-targets, the trojan gates placement and the routing metals.

2. ADDRESS TRANSLATION IN RISC-V



The targeted CPU design is **Ariane/CVA6**, a **64-bit** implementation of the **RISC-V** Instruction Set Architecture (ISA).

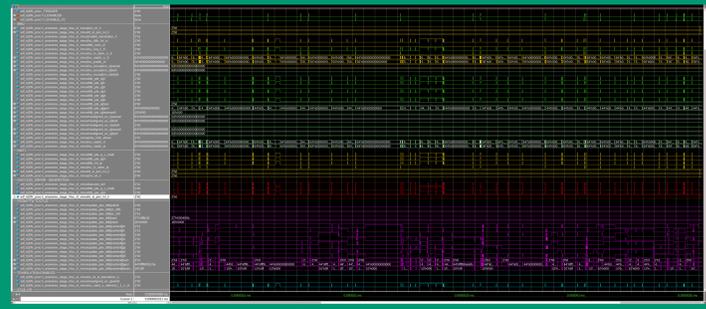


The HT targets memory protected areas of the CPU through the store memory-access mechanism.

Specific memory access checks performed during the address translation process are subverted to enable the attack.

The HT targets the **U-bit** of the **Page Table Entry (PTE)**. This bit indicates if the processor can access the page while being in **User Mode**.

Microarchitecture Level



RTL code inspection and **real-time measurements** from Ariane's internal signal states allow for behavioral profiling of the signals involved in the address translation process.

Candidate signals are then chosen for the **trigger** and **payload** circuits.

4. FUNCTIONALITY & STABILITY EVALUATION

Software Logic for HT-Testing

```

/* Kernel Programming */
#include <linux/module.h>
#include <linux/kernel.h>
#include <linux/kern_levels.h>
#include <linux/init.h>
static int vic_var;
int init_module(void)
{
    vic_var = 0xAABCCDD;
    printk(KERN_ALERT "%px = %pa\n", (void *)&vic_var, (void *)&vic_var);
    return 0;
}
void cleanup_module(void)
{
    printk(KERN_ALERT "%px = %pa\n", (void *)&vic_var, (void *)&vic_var);
    return 0;
}
MODULE_LICENSE("GPL");

/* "Malicious" User Process */
int main()
{
    u_int64_t * ptr;
    // Get the address from stdin
    printf("Memory address to be altered:");
    scanf("%p", &ptr);

    // Trigger the Trojan
    asm ("addi a5, zero, 1\n\t"
        "sll a5, a5, 0x20\n\t"
        "lui a6, 0xffff\n\t"
        "sll a6, a6, 0x15554\n\t"
        "xor a6, a6, a5\n\t"
        "lui a7, 0xffff\n\t"
        "sll a7, a7, 0x13\n\t");

    // Change the value stored in the address.
    *ptr = 0xAAAABBBB;
    return 0;
}
    
```

Once loaded by the OS, the **Linux Kernel Module (LKM)** assigns a value to a 64-bit variable defined inside the kernel address space.

The **user process** initially activates the HT and then alters the value of the variable defined by the LKM. This is not normally allowed by the OS.

Functional testing on Ariane's trojaned RTL code using the **RISC-V Assembly unit tests (Assembly Commands, Atomic Memory Operations, Floating-Point and Multiplication)**. Tests were executed error-free using Mentor's QuestaSim.

CPU **stability evaluation** through interleaved real-time execution of **Dhrystone** benchmark and **RAMSpeed** program on the FPGA implemented trojan design.

Complete **simulation** of the **Linux boot-up** process in both trojan and trojan-free Ariane. The following internal signaling comparison proved that trojan insertion leaves all the internal signals unaffected.

7. ONGOING RESEARCH

- Insertion of Astrahan in Ariane's finalized layout using ASIC tools, to showcase the HT's scalability, feasibility and low-footprint design.
- Evaluation of different HT detection methods efficacy with Astrahan's FPGA implementation.
- Further research regarding the identification of points of interest in finalized complex layouts for trojan insertion. Astrahan will be used as a prime example.

ACKNOWLEDGEMENTS

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This is an Open-Source project. Code and examples can be found in: <https://github.com/0ena/riscv-hw-trojans>



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